Theory of n-composite naturals and determination of the generating set for primes of the form 24m+1

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In this paper, concept of 'n-composite natural' has been introduced for naturals of the form 6k+1. A theory has been developed which helps in identifying naturals of the form 6k+1 for different nature of the values of k. It reveals that a natural of the form 24m+1 can be 4n-composite only which finally helps to specify the set of naturals m for which 24m+1 is prime.

Keywords: Natural numbers, Composite, Prime.

1. Introduction

To establish an effective algorithm for determining primes is of great interest in the study of number theory. There are many characterizations [1], [3], [4] and algorithms [2], [5], [6] already appeared in literature for determining primes, but the theory developed in this article will supply information on the composites of the form 6k+1, and thereafter answer the following question:

'Can we determine the particular subset of N consisting of naturals m for which 24m + 1 is prime?'

Since primes (>3) are of the form 6k - 1 or 6k+1, in this article we shall consider naturals of the form 6k+1 only and develop a theory on composite naturals of this form. This theory observes special features of composite naturals of the form 6k + 1 depending upon nature of k.

Since every composite of the form 6k + 1 has proper factors of the form 6t+1 or 6t-1 only, it clearly follows that for a composite 6k+1, 6k+1 = AB where A and B are proper factors of 6k+1 where B-A = 6n for some natural n (ignoring the case that 6k+1 is a perfect square).

Thus we define

Definition 1.1 Let the composite 6k + 1 be not a perfect square. Then 6k + 1 is said to be n-composite for some natural n if 6k+1 = AB for some proper factors A, B of 6k+1 where B-A = 6n.

For example, 55 = 6.9 + 1 and 55 = 5.11, where 11-5=6=6.1. Hence 55 is 1-composite.

Also an *n*-composite 6k+1 may be *m*-composite for some $m(\neq n)$. For example, 385 = 5.7.11. Hence it is 4-composite, 8-composite and 12-composite.

2 Properties of n-composite naturals and determination of the generating set for primes of the form 24m+1

Theorem 2.1. If 6k+1 is *n*-composite then $n \le \frac{k-4}{5}$.

Proof. Clearly A \geq 5, B \geq 5. Hence 6k +1 \geq 5(6*n* + 5). This completes the proof.

Theorem 2.2. If 6k + 1 is 2n-composite then $n \le \frac{k-4}{10}$. Proof is easy.

Theorem 2.3. If 6k + 1 is 4n-composite then $n \le \frac{k-4}{20}$.

Proof is easy.

Theorem 2.4. If 6k + 1 is *n*-composite then $9n^2 + 6k + 1$ is a perfect square.

Proof. Let 6k + 1 = AB, B - A = 6n. Then A(A + 6n) - (6k + 1) = 0 and hence $A = -3n + \sqrt{9n^2 + 6k + 1}$. If A exists then the result follows.

The following lemma will help in proving Theorem 2.6 and Theorem 2.7.

Lemma 2.5. If t is an odd natural then the equation $x^2 - y^2 = 6t$ has no positive integer solutions for x and y.

Proof. If possible, let $x^2 - y^2 = 6t$ for some naturals *x* and *y* where *t* is an odd natural.

Then (x + y)(x - y) = 6t implies that x and y cannot be even or odd simultaneously, because in that case 4 is a factor of (x + y)(x - y) but 4 is not a factor of 6t. Hence if x is even then y is odd or if x is odd then y is even. But, in any case (x + y)(x - y) is odd whereas 6t is always even. Hence a contradiction arises. This proves the lemma. **Theorem 2.6.** If 6k + 1 is n-composite then (6k + 1) + 6t is not n-composite for every odd natural t.

Proof. Let 6k + 1 be *n*-composite and if possible, let (6k + 1) + 6t be *n*-composite for some odd *t*. Let 6k + 1 + 6 = AB where B - A = 6n. Then $A = -3n + \sqrt{9n^2 + 6k + 1 + 6t}$ and since A exists, $9n^2 + 6k + 1 + 6t = y^2$ for some $y \in N$.

Then $9n^2 + 6k + 1 = y^2 - 6t$ (1)

Also by theorem 4, for some $x \in N$ we have

$$9n^2 + 6k + 1 = x^2 \tag{2}$$

Then by (1) and (2), $y^2 - x^2 = 6t$. But, by Lemma 1, this is not possible. Hence the theorem is proved.

Theorem 2.7.

(i) If k is even then 6k + 1 is not n-composite for every odd n.

(ii) If k is odd then 6k + 1 is not n-composite for every even n.

Proof. (i): Let k be even. Note that for any n, (6k + 1)(6k + 1 + 6n) is *n*-composite. Now

(6k + 1)(6k + 1 + 6n) = (6k + 1) + 6(6k + 1)(k + n) =(6k + 1) + 6t say, where t = (6k + 1)(k + n). Hence if *n* is odd then *t* is odd, where (6k + 1) + 6t is *n*-composite. So by theorem 5 it follows that 6k + 1 is not *n*-composite. Thus if *k* is even then 6k + 1 is not *n*-composite for every odd *n*.

(ii) : Let *k* be odd. Following lines of proof given in (i) we can conclude the result.

Theorem 2.9 and Theorem 2.10 are consequences of the following lemma.

Lemma 2.8. If t is any odd natural then the equation $x^2 - y^2 = 12t$ does not possess any positive odd integer solutions for x and y.

Proof. If possible, let $x^2 - y^2 = 12t$ for some odd naturals x and y where t is an odd natural.

Let $x = 2t_1 - 1$ and $y = 2t_2 - 1$. Then (x + y)(x - y) = 12timplies $(t_1 + t_2 - 1)(t_1 - t_2) = 3t$.

Since t is odd, both $t_1 + t_2 - 1$ and $t_1 - t_2$ are odd. Let $t_1 + t_2 - 1 = 2A - 1$ and $t_1 - t_2 = 2B - 1$. Adding we get

 $2t_1 = 2(A + B) - 1$, which is absurd, since L.H.S. is even whereas R.H.S. is odd. This proves the lemma.

Theorem 2.9. If 6k + 1 is 2n-composite then (6k + 1) + 12t is not 2n-composite for every odd natural t.

Proof. Let 6k + 1 be 2*n*-composite and if possible, let (6k + 1) + 12t be 2*n*-composite for some odd *t*. Let 6k + 1 + 12t = AB where B - A = 6.2n = 12n. Then $A = -6n + \sqrt{36n^2 + 6k + 1 + 12t}$ and since *A* exists, $36n^2 + 6k + 1 + 12t = y^2$ for some $y \in N$.

Then $36n^2 + 6k + 1 = \gamma^2 - 12t$ (3)

Also by theorem 4, for some $x \in N$ we have

 $9.(2n)^2 + 6k + 1 = x^2 \tag{4}$

Then by (3) and (4), $y^2 - x^2 = 12t$. Note that x and y are both odd. But, by lemma 2, this is not possible. Hence the theorem is proved.

Theorem 2.10. If k = 4m for any natural m then 6k + 1 is not composite for every odd n.

Proof. Note that for any n, (6k + 1)(6k + 1 + 12n) is 2*n*-composite. Now

(6k + 1)(6k + 1 + 12n) = (6k + 1) + 6(6k + 1)(k + 2n)= (6k + 1) + 12t say, where $t = (6k + 1)\frac{k + 2n}{2}$. Now k = 4m. So t = (6k + 1)(2m + n). Hence if *n* is odd then *t* is odd, where (6k + 1) + 12t is 2*n*-composite. So by theorem 7 it follows that 6k + 1 is not 2*n*-composite. Thus if k = 4m then 6k + 1 is not 2*n*-composite for every odd *n*.

Now we have the following result.

Corollary 2.11. If k = 4m then

(i) 6k + 1 is not n-composite for every odd n and

(ii) 6k + 1 is not 2n-composite for every odd n.

Proof follows from theorem 2.7 and theorem 2.11.

As a consequence we get

Theorem 2.12. Let k = 4m and 6k + 1 be not a perfect square. Then the following statements are equivalent :

(i) 24m + 1 is prime.

(ii) $24m + 1 \neq t(24n + t)$ for all $n \leq \frac{m-1}{5}$ and any t of the form 6v - 1 or 6v + 1.

Proof. (ii) \rightarrow (*i*). Let 24m + 1 be composite. Then by corollary 2.11, it can only be 4n-composite. Hence by

theorem 2.3, $n \le \frac{m-1}{5}$. Now by theorem 2.4, we have $9(4n)^2 + 24m + 1 = y^2$ for some natural y. Then

 $9(4n)^2 + 24m + 1 = y^2 = (12n + t)^2$ (5)

for some natural *t*. So 24m + 1 = t(24n + t). Now L.H.S. of (5) is of the form 6w + 1. Hence *t* must be of the form 6v - 1 or 6v + 1. This completes the proof.

(i) \rightarrow (ii). This is obvious.

Corollary 2.13. Let k = 4m and 6k + 1 be not a perfect square. Then the following statements are equivalent :

(i) 24m + 1 is prime.

(ii)
$$m \neq nt + \frac{t^2 - 1}{24}$$
 for all $n \leq \frac{m - 1}{5}$ and any t of the

form 6v - 1 or 6v + 1.

(iii) 24m + 1 is not congruent to $t^2 \pmod{24t}$, for any $t \left(< \sqrt{24m + 1} \right)$ of the form 6v - 1 or 6v + 1.

Proof is evident.

Now we consider the following important question:

Can we completely determine the particular subset P(24) (say) of N such that for each $m \in P(24)$, 24m + 1 is prime and conversely if 24m + 1 is prime then $m \in P(24)$?

The answer to this question is as follows :

Consider the following subsets of *N*. $TF(1) = \{nt + \frac{t^2 - 1}{24} : n \in N, t \in N, t = 6v - 1 \text{ or } t = 6v + 1, v \in N\}.$

 $TF(2) = \{m : m \in N, 24m + 1 = y^2, \text{ for some } y \in N\}.$

Members of *TF*(1) are of the forms, 5n + 1, 7n + 2, 11n + 5, 13n + 7, 17n + 12,..... and so on.

Let $S = TF(1) \cup TF(2)$.

Then it follows by corollary 2.13 that the subset P(24) = N - S of N will be the Prime determining subset for the naturals of the form 24m + 1. This means, for every $m \in P(24)$, 24m + 1 is prime and conversely, if 24m + 1 is prime then $m \in P(24)$.

Writing members of S in ascending order we can enumerate members of P(24) in ascending order.

One can verify the main result obtained in this article for any finite subset of N by using the following algorithm.

Algorithm :

1. Set limit. 2. Generate set TF(2): 2.1. Set $TF(2) \leftarrow \{\}$. 2.2. Set $m \leftarrow 1$. 2.3. While $m \leq \text{limit}$ 2.3.1. Set $\gamma^2 + 24m + 1$ 2.3.2. If y^2 is a perfect square, add *m* to TF(2). 2.3.3. Set $m \leftarrow m + 1$ and iterate loop. 3. Generate set TF(1): 3.1. Set $S_1 \leftarrow \{\}, S_2 \leftarrow \{\}, \text{flag} \leftarrow 0$. 3.2. Set $V \leftarrow 1$. 3.3. While $V \leq \text{limit}$ 3.3.1. If flag =1, break loop. 3.3.2. Set $n \leftarrow 1$. 3.3.3. While $n \leq limit$ 3.3.3.1. Set $t \leftarrow 6V - 1$, $t_2 \leftarrow 6V + 1$. 3.3.3.2. Set $a \leftarrow n.t_1 + \frac{1}{24}(t_1^2 - 1)$. 3.3.3.3. If $a \leq \text{limit}$, add a to S_1 , else if n = 1, set flag \leftarrow 1 and break current loop, else break current loop. 3.3.3.4. Set $b \leftarrow n.t_2 + \frac{1}{24}(t_2^2 - 1)$. 3.3.3.5. If $b \leq limit$, add b to S_2 . 3.3.3.6. Set $n \leftarrow n + 1$ and iterate loop. 3.3.4. Set $V \leftarrow V + 1$ and iterate loop. 3.4. Set $TF(1) + S_1 \cup S_2$. 4. Set $S \leftarrow TF(1) \cup TF(2)$. 5. Set $N \leftarrow \{1, 2, 3, \dots, (\text{largest element of } S)\}$. 6. Set $P(24) \leftarrow N-S$. 7. Set $P \leftarrow \{24p + 1 : p \in P(24)\}.$

8. If all elements of *P* are prime, declare success, else declare failure.

Conclusion :

It is worthwhile to mention that primes are important in Computer Science and Cryptography because the security of many encryption algorithms is based on the fact that the multiplication of two large prime numbers is very fast, but takes a lot of processing to do the reverse.

This article develops a method which gives us the particular values of m for which 24m + 1 are prime numbers. An algorithm has been constructed which helps in finding larger values of m and the corresponding values of the primes 24m + 1. In this way, this algorithm will certainly be helpful in obtaining larger prime numbers which will be useful in developing certain encryption techniques.

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